

**Note**

**A Method and Two Algorithms on the Theory of Partitions**

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In this note we describe a general principle for selecting at random from a collection of combinatorial objects, where "at random" means in such a way that each of the objects has equal probability, *a priori*, of being selected. We apply this principle by displaying two algorithms, the first of which will select a random partition of an integer  $n$ , and the second will choose a random partition of a set  $S$  of  $n$  elements.

We begin by stating the general principle, albeit rather vaguely. Let  $a_n$  be the number of combinatorial objects of order  $n$ . Suppose the numbers  $a_n$  satisfy a recurrence relation of the form

$$a_n = \sum_{m < n} \alpha_{m,n} a_m, \tag{1}$$

in which  $\alpha_{m,n} \geq 0$ . Suppose further that the recurrence (1) can be given a combinatorial interpretation, by which we mean that there is a proof of (1) in which, by an explicit construction, the  $a_m$  objects of order  $m$  are extended to  $\alpha_{m,n} a_m$  objects of order  $n$ . Then we have an algorithm for selecting an object of order  $n$  at random: first choose a value of  $m$ ,  $0 \leq m \leq n - 1$  according to the probabilities

$$\text{Prob}(m) = \frac{\alpha_{m,n} a_m}{a_n} \quad (m = 0, 1, \dots, n - 1). \tag{2}$$

Then, inductively, select a random object of order  $m$ , and extend it as described in the proof.

We make the above statement precise in two cases. First, if  $p(n)$  is the

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number of partitions of the integer  $n$ , then an identity of Euler asserts that

$$np(n) = \sum_{m < n} \sigma(n - m) p(m) \quad (p(0) = 1), \quad (3)$$

where  $\sigma(m)$  is the sum of the divisors of the integer  $m$ , while if  $a_n$  is the number of partitions of a *set* of  $n$  elements, then the  $a_n$  satisfy

$$a_n = \sum_{m=0}^{n-1} \binom{n-1}{m} a_m \quad (a_0 = 1). \quad (4)$$

Next we must ask for the combinatorial interpretations of (3), (4). For (3), let  $\pi$  denote a fixed partition of an integer  $m < n$ . Let  $d$  be a divisor of  $n - m$ . With the pair  $(\pi, d)$  we associate exactly  $d$  copies of a single partition  $\pi'$ , of  $n$ . Here  $\pi'$  is obtained by adjoining to the partition  $\pi$ , of  $m$ , exactly  $(n - m)/d$  copies of  $d$ .

We claim that as  $d$  varies over all divisors of  $n - m$  and  $\pi$  varies over all partitions of  $m$  ( $m = 0, 1, \dots, n - 1$ ), each partition of  $n$  is constructed exactly  $n$  times by this process, which will establish (3). Indeed, let

$$\pi' : n = \mu_1 r_1 + \mu_2 r_2 + \dots + \mu_k r_k \quad (5)$$

be a fixed partition of  $n$ , where the  $r_i$  are the *distinct* parts of  $\pi'$  and the  $\mu_i$  are their multiplicities. Then  $\pi'$  is constructed by adjoining  $t$  copies of  $r_i$  to a partition of  $n - tr_i$  for each  $1 \leq t \leq \mu_i$ ,  $1 \leq i \leq k$ , and by replicating the resulting partition of  $n$   $r_i$  times, which gives a total of

$$\sum_{i=1}^k r_i \sum_{t=1}^{\mu_i} 1 = \mu_1 r_1 + \dots + \mu_k r_k = n,$$

times altogether, as claimed.

As for partitions of sets, the interpretation of (4) is well known, but we include it for completeness: If  $\pi$  is a partition of  $\{1, 2, \dots, m\}$ , let  $S$  be a subset  $\{s_1, \dots, s_m\}$  of  $m$  elements chosen from  $\{1, 2, \dots, n - 1\}$ . With  $(\pi, S)$  we associate the partition of  $\{1, 2, \dots, n\}$  in which elements  $s_i$  and  $s_j$  from  $S$  are in the same class iff their subscripts  $i$  and  $j$  are in the same class of  $\pi$ , while all elements not in  $S$  are in a single class. It is easy to check that each partition of  $\{1, 2, \dots, n\}$  occurs exactly once. (We have assumed  $s_1 < s_2 < \dots < s_m$ .)

We can now describe the algorithms.

**ALGORITHM I.** *Given  $n$ , select a random partition of  $n$ .*

- (A) Set  $n' \leftarrow n$ ,  $P \leftarrow$  empty partition.
- (B) Choose an integer  $m < n'$  according to the probabilities

$$\text{Prob}(m) = \frac{\sigma(n' - m) p(m)}{n' p(n')} \quad (m = 0, 1, \dots, n' - 1).$$

- (C) Choose a divisor  $d$  of  $n' - m$  according to the probabilities

$$\text{Prob}(d) = d/\sigma(n' - m) \quad (d \mid (n' - m)).$$

- (D) Adjoin to the partition  $P$   $(n' - m)/d$  copies of  $d$ .
- (E) Replace  $n'$  by  $m$ .
- (F) If  $n' = 0$ , stop. Otherwise return to step (B).

ALGORITHM S. Given  $n$ , select a random partition of  $U = \{1, 2, \dots, n\}$ .

- (A) Set  $U' \leftarrow U$ ,  $m \leftarrow n$ ,  $P \leftarrow$  unique partition of empty set.
- (B) Choose an integer  $k < m$  according to the probabilities

$$\text{Prob}(k) = \binom{m-1}{k} \frac{a_k}{a_n} \quad (0 \leq k \leq m-1).$$

- (C) Let  $l$  be the largest element of  $U'$ . Choose a random  $k$ -subset  $S$  of  $U' - \{l\}$ .

- (D) Adjoin  $U' - S$  as a single class to  $P$ . Stop if  $k = 0$ . Otherwise set  $m \leftarrow k$ ,  $U' \leftarrow S$  and return to step (B).

It remains to show that all partitions have equal probability ( $= p(n)^{-1}$  or  $a_n^{-1}$ , respectively). We show this in the case of partitions of an integer, the other case being similar.

Let

$$\pi' : n = \mu_1 r_1 + \dots + \mu_k r_k \tag{6}$$

be a fixed partition of  $n$ . Inductively, suppose that for all  $n' < n$  it has been shown that our algorithm produces partitions of  $n'$  with all equal probabilities. Then  $\text{Prob}(\pi')$  is a sum over all partitions  $\pi''$  of integers  $m < n$  of  $\text{Prob}(\pi'')$  multiplied by the probability that on the next step  $\pi''$  is extended to  $\pi'$ :

$$\begin{aligned} \text{Prob}(\pi') &= \sum \text{Prob}(\pi'') \text{Prob}(\pi''_m \rightarrow \pi') \\ &= \sum 1/p(m) \text{Prob}(\pi''_m \rightarrow \pi'). \end{aligned} \tag{7}$$

The last factor vanishes unless  $\pi'$  can result from  $\pi''_m$  by adjunction of

exactly  $t$  copies of one part  $r_i$ . In the latter case, if  $\pi'$  is the partition (6) then

$$\begin{aligned} \text{Prob}(\pi_m'' \rightarrow \pi') &= \text{Prob}(m = n - tr_i) \text{Prob}(d = r_i) \\ &= \left\{ \frac{\sigma(tr_i) p(m)}{np(n)} \right\} \left\{ \frac{r_i}{\sigma(tr_i)} \right\} \\ &= \frac{p(m)}{np(n)} r_i. \end{aligned}$$

The sum (7) over all candidates  $\pi_m''$  for extension to  $\pi'$  in a single step becomes just

$$\begin{aligned} \text{Prob}(\pi') &= \sum_{i=1}^k \sum_{t=1}^{\mu_i} \frac{1}{p(m)} \frac{p(m)}{n p(n)} r_i \\ &= \frac{1}{n p(n)} \sum_{i=1}^k \mu_i r_i \\ &= \frac{1}{p(n)}, \end{aligned}$$

as required.

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*Note added in proof.* A generalization of these ideas, along with other applications, appears in [1].

#### REFERENCE

1. A. NIJENHUIS AND H. S. WILF, "Combinatorial Algorithms," Academic Press, New York, 1975.